



Recursion (cont'ed)

Tecniche di Programmazione – A.A. 2014/2015



Summary

1. Definition and divide-and-conquer strategies
2. Simple recursive algorithms
 1. Fibonacci numbers
 2. Dicothomic search
 3. X-Expansion
 4. Proposed exercises
3. Recursive vs Iterative strategies
4. More complex examples of recursive algorithms
 1. Knight's Tour
 2. Proposed exercises

Recursion

► Divide et Impera

- ▶ Split a problem \mathcal{P} into $\{\mathcal{Z}_i\}$ where \mathcal{Z}_i are still complex, yet *simpler* instances of the same problem.
- ▶ Solve $\{\mathcal{Z}_i\}$, then merge the solutions
- ▶ Merge & split must be “simple”
- ▶ A.k.a., *Divide n' Conquer*

► Exploration

- ▶ Systematic procedure to enumerate all possible solutions
- ▶ Solutions \leftrightarrow Paths
- ▶ Similar to D+I with $\{\mathcal{A}, \mathcal{P}'\}$

► 3

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Complessità



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Divide et Impera

```

▶ Solution = Solve ( Problem ) ;

▶ Solve ( Problem ) {
    ▶ List<SubProblem> subProblems = Divide ( Problem ) ;
    ▶ For ( each subP[i] in subProblems ) {
        ▶ SubSolution[i] = Solve ( subP[i] ) ;
    }
    ▶ Solution = Combine ( SubSolution[ 1..N ] ) ;
    ▶ return Solution ;
}

```

▶ 5

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Divide et Impera – Divide and Conquer

```

▶ Solution = Solve ( Problem ) ;

▶ Solve ( Problem ) {
    ▶ List<SubProblem> subProblems = Divide ( Problem ) ;
    ▶ For ( each subP[i] in subProblems ) {
        ▶ SubSolution[i] = Solve ( subP[i] ) ;
    }
    ▶ Solution = Combine ( SubSolution[ 1..N ] ) ;
    ▶ return Solution ;
}

```

▶ 6

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“a” sub-problems, each
“b” times smaller than
the initial problem

recursive call

Divide et Impera – Divide and Conquer

```

▶ Solve ( Problem ) {
    ▶ if( problem is trivial )
        ▶ Solution = Solve_trivial ( Problem ) ;
    ▶ else {
        ▶ List<SubProblem> subProblems = Divide ( Problem ) ;
        ▶ For ( each subP[i] in subProblems ) {
            □ SubSolution[i] = Solve ( subP[i] ) ;
        ▶ }
        ▶ Solution = Combine ( SubSolution[ 1..N ] ) ;
    ▶ }
    ▶ return Solution ;
}

```

do recursion

▶ 7

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What about complexity?

- ▶ a = number of sub-problems for a problem
- ▶ b = how smaller sub-problems are than the original one
- ▶ n = size of the original problem
- ▶ $T(n)$ = complexity of **Solve**
 - ▶ ...our unknown complexity function
- ▶ $\Theta(1)$ = complexity of **Solve_trivial**
 - ▶ ...otherwise it wouldn't be trivial
- ▶ $D(n)$ = complexity of **Divide**
- ▶ $C(n)$ = complexity of **Combine**

▶ 8

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What about complexity?

- ▶ **$\Theta(1)$** = complexity of **Solve_trivial**
 - ▶ ...otherwise it wouldn't be trivial
- ▶ **D(n)** = complexity of **Divide**
- ▶ **C(n)** = complexity of **Combine**

- ▶ **$T(n) = \Theta(n \log n)$**

▶ 9

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Exploration

- ▶ **Explore () {**
 - ▶ List<Step> steps = **PossibleSteps** (Problem) ;
 - ▶ For (each s in steps) {
 - ▶ **Do** (s)
 - ▶ **Explore** () ;
 - ▶ **Undo** (s)
 - ▶ }
- ▶ }

▶ 10

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Exploration

```

▶ Explore () {
    ▶ List<Step> steps = PossibleSteps ( Problem ) ;
    ▶ For ( each s in steps ) {
        ▶ Do ( s )
        ▶ Explore () ;           Update “global” status
        ▶ Undo ( s )
    }
}

```

▶ 11

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What about complexity?

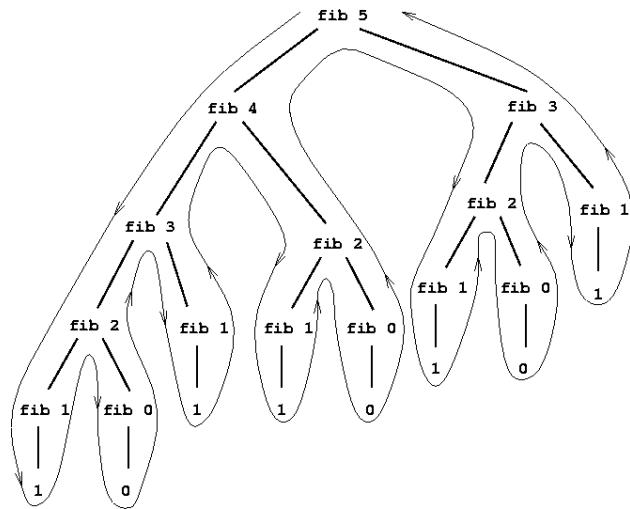
- ▶ (Almost always)

- ▶ $T(n) = \Theta(e^n)$

▶ 12

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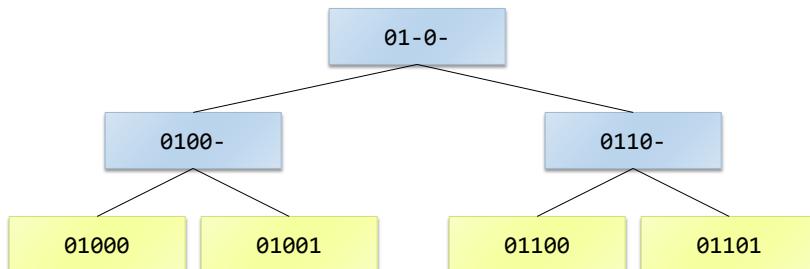
Recursion Tree (exploration)



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Don't care expansion



▶ 14

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Recursive vs Iterative strategies

Recursion

Recursion and iteration

- ▶ Every **recursive** program can **always** be implemented in an **iterative** manner
- ▶ The best solution, in terms of efficiency and code clarity, depends on the problem



More complex examples of recursive algorithms

Recursion

The Knapsack Problem

Input: Weight of N items $\{w_1, w_2, \dots, w_n\}$
Cost of N items $\{c_1, c_2, \dots, c_n\}$
Knapsack limit S

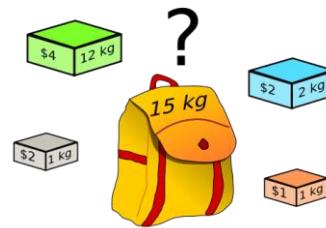
Output: Selection for knapsack: $\{x_1, x_2, \dots, x_n\}$ where $x_i \in \{0, 1\}$.

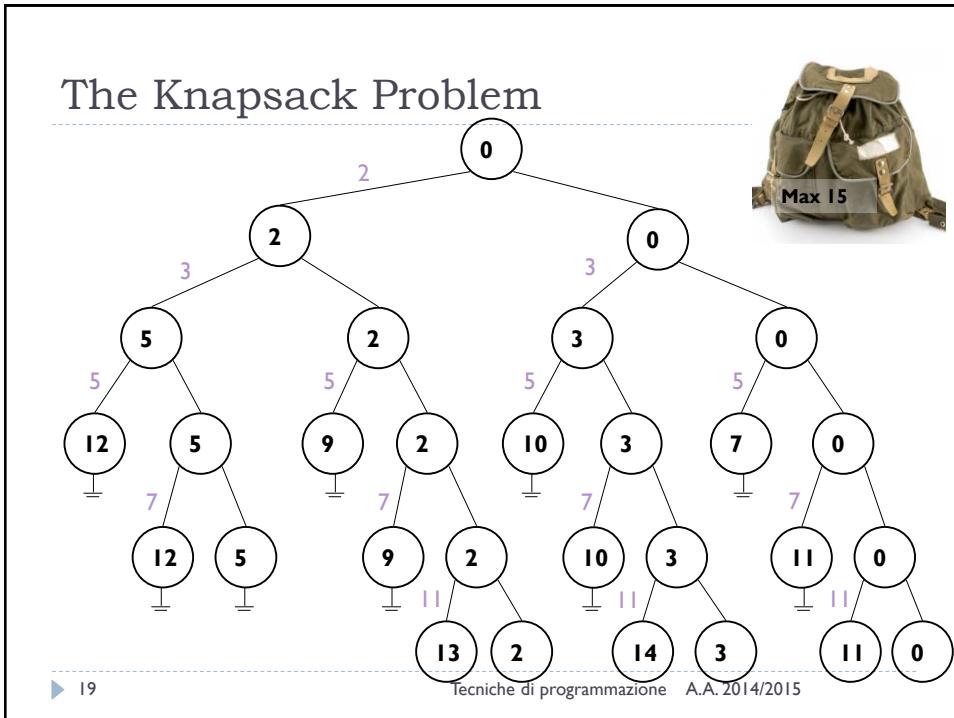
Sample input:

$$w_i = \{1, 1, 2, 4, 12\}$$

$$c_i = \{1, 2, 2, 10, 4\}$$

S=15





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